

## LECTURE 23

In this lecture, we define  $\mathcal{O}$ -modules for an  $\infty$ -operad  $\mathcal{O}^\otimes$ .

### 1. OUTLINE

1.1. In the previous lectures, we constructed operads  $\mathbf{CMod}^\otimes$ ,  $\mathbf{LMod}^\otimes$ ,  $\mathbf{RMod}^\otimes$  and  $\mathbf{BMod}^\otimes$  such that the corresponding *algebra objects* in a symmetric monoidal  $\infty$ -category are pairs/triples of commutative/associative algebras and certain types of modules of them. However, these constructions are conducted in an *ad hoc* manner based on the explicit description of  $\mathbf{Comm}^\otimes$  and  $\mathbf{Assoc}^\otimes$ . In this lecture, we give a general theory of modules for  $\mathcal{O}$ -algebras for *coherent*  $\infty$ -operad  $\mathcal{O}^\otimes$ .

1.2. The definition of coherent  $\infty$ -operads will be given later. For now, let us treat it as a blackbox.

Let  $\mathcal{O}^\otimes$  be a coherent  $\infty$ -operad and  $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$  be an  $\mathcal{O}$ -monoidal  $\infty$ -category. Let  $A \in \mathbf{Alg}_{/\mathcal{O}}(\mathcal{C})$  be an  $\mathcal{O}$ -algebra. For each color  $X \in \mathcal{O}$ , we will construct an  $\infty$ -category,  $\mathbf{Mod}_A^{\mathcal{O}}(\mathcal{C})_X^\otimes$ , called the  $\infty$ -category of  $\mathcal{O}$ - $A$ -modules in  $\mathcal{C}$  of color  $X$ . There will be a forgetful functor

$$\mathbf{Mod}_A^{\mathcal{O}}(\mathcal{C})_X^\otimes \rightarrow \mathcal{C}_X^\otimes$$

which sends an  $\mathcal{O}$ - $A$ -module to its **underlying object**. In fact, we will allow  $X$  to be a *multi-color*, i.e., any object in  $\mathcal{O}^\otimes$ . Moreover, we will have a *parameterized* version

$$(1.1) \quad \mathbf{Mod}^{\mathcal{O}}(\mathcal{C})^\otimes \rightarrow \mathbf{Alg}_{/\mathcal{O}}(\mathcal{C}) \times \mathcal{O}^\otimes$$

such that its fiber at  $(A, X)$  is  $\mathbf{Mod}_A^{\mathcal{O}}(\mathcal{C})_X^\otimes$ . Under *good conditions*, we have the following results, details of which will be explained in the next lecture.

- For fixed  $A \in \mathbf{Alg}_{/\mathcal{O}}(\mathcal{C})$ , the fiber of (1.1)

$$\mathbf{Mod}_A^{\mathcal{O}}(\mathcal{C})^\otimes \rightarrow \mathcal{O}^\otimes$$

will be an  $\mathcal{O}$ -monoidal  $\infty$ -category. This generalizes the relative tensor products  $-\otimes_A-$  in commutative algebra.

- The projection

$$\mathbf{Mod}^{\mathcal{O}}(\mathcal{C})^\otimes \rightarrow \mathbf{Alg}_{/\mathcal{O}}(\mathcal{C})$$

will be a Cartesian fibration. This generalizes the restriction functors in commutative algebra.

- Note however that for a morphism  $A \rightarrow B$  in  $\mathbf{Alg}_{/\mathcal{O}}(\mathcal{C})$ , the contravariant transport functor

$$\mathbf{res}: \mathbf{Mod}_B^{\mathcal{O}}(\mathcal{C})^\otimes \rightarrow \mathbf{Mod}_A^{\mathcal{O}}(\mathcal{C})^\otimes$$

is *not* an  $\mathcal{O}$ -monoidal functor. Instead, it is only a *right lax*  $\mathcal{O}$ -monoidal one, i.e., an  $\infty$ -operad map. This generalizes the corresponding fact in commutative algebra.

- For fixed  $X \in \mathcal{O}^\otimes$ , the fiber of (1.1)

$$\mathrm{Mod}^\mathcal{O}(\mathbb{C})_X^\otimes \rightarrow \mathrm{Alg}_{/\mathcal{O}}(\mathbb{C})$$

will be a coCartesian fibration. This generalizes the induction functors in commutative algebra.

- Note however that the functors

$$\mathrm{ind} : \mathrm{Mod}_A^\mathcal{O}(\mathbb{C})_X^\otimes \rightarrow \mathrm{Mod}_B^\mathcal{O}(\mathbb{C})_X^\otimes, X \in \mathcal{O}^\otimes$$

cannot be assembled to a functor over  $\mathcal{O}^\otimes$ . Indeed, as the left adjoint of the right lax  $\mathcal{O}$ -monoidal functor  $\mathrm{res}$ , the functor  $\mathrm{ind}$  should in general only be *left lax*  $\mathcal{O}$ -monoidal. In the commutative case, the induction functor happens to be  $\mathcal{O}$ -monoidal, but this already fails for  $\mathcal{O}^\otimes = \mathrm{Assoc}^\otimes$ . See Exercise 1.4 below.

**Warning 1.3.** For  $\mathcal{O}^\otimes = \mathrm{Assoc}^\otimes$ , objects in  $\mathrm{Mod}_A^{\mathrm{Assoc}^\otimes}(\mathbb{C})$  are  $(A, A)$ -bimodules rather than left or right  $A$ -modules. Indeed, to define relative tensor products over  $A$ , we need a left  $A$ -module and a right  $A$ -module.

**Exercise 1.4.** Let  $A \rightarrow B$  be a homomorphism between associative rings. Show that the restriction functor

$$\mathrm{res} : {}_B\mathrm{BMod}_B \rightarrow {}_A\mathrm{BMod}_A$$

has a left adjoint

$$\mathrm{ind} : {}_A\mathrm{BMod}_A \rightarrow {}_B\mathrm{BMod}_B$$

which is left lax monoidal but not monoidal.

## 2. COHERENT $\infty$ -OPERADS

2.1. The  $\infty$ -category  $\mathrm{Mod}^\mathcal{O}(\mathbb{C})_X^\otimes$  will be defined as a certain full sub- $\infty$ -category

$$\mathrm{Mod}^\mathcal{O}(\mathbb{C})_X^\otimes \subseteq \mathrm{Fun}_{\mathcal{O}^\otimes}(\mathbb{K}_X, \mathbb{C}^\otimes),$$

where  $\mathbb{K}_X$  is a certain  $\infty$ -category over  $\mathcal{O}^\otimes$  that plays the role of  $\mathrm{CMod}^\otimes$  in the  $\mathrm{Comm}$ -case.

We will obtain the functoriality of  $\mathrm{Mod}^\mathcal{O}(\mathbb{C})_X^\otimes$  in  $X$  via that of  $\mathbb{K}_X$ . In other words, we will define a certain  $\infty$ -category  $\mathbb{K}_\mathcal{O}$  equipped with *two* projections

$$(\mathrm{ev}_0, \mathrm{ev}_1) : \mathbb{K}_\mathcal{O} \rightarrow \mathcal{O}^\otimes \times \mathcal{O}^\otimes$$

and define

$$\mathbb{K}_X := \{X\} \times_{\mathcal{O}^\otimes, \mathrm{ev}_0} \mathbb{K}_\mathcal{O} \xrightarrow{\mathrm{ev}_1} \mathcal{O}^\otimes$$

where the fiber product uses  $\mathrm{ev}_0$ , and the projection  $\mathbb{K}_X \rightarrow \mathcal{O}^\otimes$  uses  $\mathrm{ev}_1$ .

To achieve this, we will *not* make  $\mathbb{K}_X$  an  $\infty$ -operad. When  $\mathcal{O}^\otimes = \mathrm{Comm}^\otimes$  and  $X = \langle 1 \rangle$ , this means we only allow objects in  $\mathrm{CMod}^\otimes$  with at most one  $\mathfrak{m}$ -term.

2.2. To define  $\mathbb{K}_\mathcal{O}$ , we need some terminologies.

**Definition 2.3.** Let  $\mathcal{O}^\otimes$  be an  $\infty$ -operad. We say a morphism  $f : (X_1, \dots, X_m) \rightarrow (Y_1, \dots, Y_n)$  lying over  $\langle m \rangle \xrightarrow{\alpha} \langle n \rangle$  is **semi-inert** if for each  $j \in \langle n \rangle^\circ$ , exactly one of the following conditions holds:

- The set  $\alpha^{-1}(\{j\})$  is empty.
- The set  $\alpha^{-1}(\{j\})$  contains exactly one element  $i$ , and the corresponding morphism  $X_i \rightarrow Y_j$  is an isomorphism.

**Exercise 2.4.** Show that  $f$  is inert iff (b) holds for each  $j \in \langle n \rangle^\circ$ .

**Definition 2.5.** Let

$$\mathbf{K}_O \subseteq \text{Fun}([1], \mathbf{O}^\otimes)$$

be the full sub- $\infty$ -category of semi-inert morphisms in  $\mathbf{O}^\otimes$  and

$$\mathbf{K}_X := \{X\}_{\mathbf{O}^\otimes, \text{ev}_0} \times \mathbf{K}_O$$

be the  $\infty$ -category of semi-inert morphisms out of  $X$ .

**Definition 2.6.** We say a morphism in  $\mathbf{K}_O$  is *inert* iff both functors  $\text{ev}_0$  and  $\text{ev}_1$  send it to an inert morphism in  $\mathbf{O}^\otimes$ .

**Exercise 2.7.** For  $\mathbf{O}^\otimes = \text{Comm}^\otimes$  and  $X = \langle 1 \rangle$ , describe the category  $\mathbf{K}_X$  and construct a canonical functor  $\mathbf{K}_X \rightarrow \text{CMod}^\otimes$  defined over  $\text{Comm}^\otimes$ . Show that

$$\text{Alg}_{\text{CMod}}(\mathbb{C}) \rightarrow \text{Fun}_{\text{Comm}^\otimes}(\text{CMod}^\otimes, \mathbb{C}^\otimes) \rightarrow \text{Fun}_{\text{Comm}^\otimes}(\mathbf{K}_X, \mathbb{C}^\otimes)$$

is fully faithful and the essential image consists of those functors  $\mathbf{K}_X \rightarrow \mathbb{C}^\otimes$  that preserves inert morphisms.

2.8. We are ready to give the definition of coherent  $\infty$ -operads.

**Definition 2.9.** Let  $\mathbf{O}^\otimes$  be a small  $\infty$ -operad. We say  $\mathbf{O}^\otimes$  is *coherent* if it satisfies the following conditions

- (i) The  $\infty$ -operad  $\mathbf{O}^\otimes$  is unital.
- (ii) The  $\infty$ -category  $\mathbf{O}$  of colors is a space.
- (iii) The base-change functor along  $\text{ev}_0 : \mathbf{K}_O \rightarrow \mathbf{O}^\otimes$

$$\text{ev}_0^* : (\text{Cat}_\infty)_{/\mathbf{O}^\otimes} \rightarrow (\text{Cat}_\infty)_{/\mathbf{K}_O}, D \mapsto D \times_{\mathbf{O}^\otimes, \text{ev}_0} \mathbf{K}_O$$

admits a right adjoint  $\text{ev}_{0,*}$ .

**Proposition 2.10** (HA.3.3.1.12, 4.1.1.20). The  $\infty$ -operads  $\text{Comm}^\otimes$  and  $\text{Assoc}^\otimes$  are coherent.

2.11. Although we have not yet defined  $\mathbf{E}_k$ , let us record the following result.

**Theorem 2.12** (HA.5.1.1.1). The  $\infty$ -operad  $\mathbf{E}_k^\otimes$  is coherent for any  $k \geq 0$ .

**Remark 2.13.** Using the results in HA.B.3, one can show (iii) is equivalent to

- (iii') The functor  $\text{ev}_0 : \mathbf{K}_O \rightarrow \mathbf{O}^\otimes$  is a flat categorical fibration.

Hence by HA.3.3.2.2, our definition of coherent  $\infty$ -operads coincides with Lurie's definition in HA.3.3.1.1. In fact, under these equivalent conditions, the adjunction

$$\text{ev}_0^* : (\text{Cat}_\infty)_{/\mathbf{O}^\otimes} \rightleftarrows (\text{Cat}_\infty)_{/\mathbf{K}_O} : \text{ev}_{0,*}$$

can be realized as a Quillen adjunction between the (slice) Joyal model categories. It follows that for a categorical fibration  $D$  over  $\mathbf{K}_O$ , the functor  $\text{ev}_{0,*}$  sends it to the quasi-category  $\text{ev}_{0,*}(D)$  over  $\mathbf{O}^\otimes$  satisfying the following univresal property

$$\text{Hom}_{(\text{Set}_\Delta)_{/\mathbf{O}^\otimes}}(J, \text{ev}_{0,*}(D)) \simeq \text{Hom}_{(\text{Set}_\Delta)_{/\mathbf{K}_O}}(J \times_{\mathbf{O}^\otimes} \mathbf{K}_O, D),$$

where the fiber product in the RHS is the naive one.

**Exercise 2.14.** Let  $D \in (\text{Cat}_\infty)_{/\mathbf{O}^\otimes}$ , show that the fiber of  $\text{ev}_{0,*} \circ \text{ev}_1^*(D)$  at  $X \in \mathbf{O}^\otimes$  is equivalent to  $\text{Fun}_{\mathbf{O}^\otimes}(\mathbf{K}_X, D)$ .

## 3. DEFINITION OF MODULES

3.1. We are finally ready to construct  $\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes}$ .

**Definition 3.2.** Let  $\mathcal{O}^{\otimes}$  be a coherent  $\infty$ -operad and  $\mathcal{C}^{\otimes} \rightarrow \mathcal{O}^{\otimes}$  be an  $\infty$ -operad map. We define

$$\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes} \subseteq \text{ev}_{0,*} \circ \text{ev}_1^*(\mathcal{C}^{\otimes})$$

to be the full sub- $\infty$ -category consisting of those objects such that the corresponding functor  $\mathcal{K}_X \rightarrow \mathcal{C}^{\otimes}$  (see Exercise 2.14) preserves inert morphisms.

3.3. Note that one projection in (1.1) is automatic: we have a structure functor

$$\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes} \rightarrow \mathcal{O}^{\otimes}.$$

**Exercise 3.4.** For  $\mathcal{O}^{\otimes} = \text{Comm}^{\otimes}$ , identify the fiber of the above functor at  $\langle 1 \rangle$  with  $\text{Alg}_{\text{CMod}}(\mathcal{C})$ .

3.5. To obtain the other projection

$$\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes} \rightarrow \text{Alg}_{\mathcal{O}}(\mathcal{C}),$$

we need more definitions.

**Definition 3.6.** Let  $\mathcal{O}^{\otimes}$  be an  $\infty$ -operad. We say a morphism  $f : (X_1, \dots, X_m) \rightarrow (Y_1, \dots, Y_n)$  lying over  $\langle m \rangle \xrightarrow{\alpha} \langle n \rangle$  is **null** if for each  $j \in \langle n \rangle^{\circ}$ ,

(a) The set  $\alpha^{-1}(\{j\})$  is empty.

Let  $\mathcal{K}_{\mathcal{O}}^0 \subseteq \mathcal{K}_{\mathcal{O}}$  be the full sub- $\infty$ -category of null morphisms.

**Exercise 3.7.** Show that  $f$  is null iff it factors through the object  $\langle 0 \rangle \in \mathcal{O}^{\otimes}$ .

**Exercise 3.8.** Show that the functor  $(\text{ev}_0^0, \text{ev}_1^0) : \mathcal{K}_{\mathcal{O}}^0 \rightarrow \mathcal{O}^{\otimes} \times \mathcal{O}^{\otimes}$  is an equivalence.

**Exercise 3.9.** Show that the diagram

$$\begin{array}{ccc} & \mathcal{K}_{\mathcal{O}}^0 & \\ \text{ev}_0^0 \swarrow & \downarrow & \searrow \text{ev}_1^0 \\ \mathcal{O}^{\otimes} & \mathcal{K}_{\mathcal{O}} & \mathcal{O}^{\otimes} \\ \text{ev}_0 \longleftarrow & & \longrightarrow \text{ev}_1 \end{array}$$

induces a natural functor

$$\text{ev}_{0,*} \circ \text{ev}_1^*(\mathcal{C}^{\otimes}) \rightarrow \text{ev}_{0,*}^0 \circ \text{ev}_1^{0,*}(\mathcal{C}^{\otimes}) \simeq \mathcal{O}^{\otimes} \times \text{Fun}_{\mathcal{O}^{\otimes}}(\mathcal{O}^{\otimes}, \mathcal{C}^{\otimes})$$

which restricts to a functor

$$\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes} \rightarrow \mathcal{O}^{\otimes} \times \text{Alg}_{\mathcal{O}}(\mathcal{C})$$

defined over  $\mathcal{O}^{\otimes}$ .

3.10. Let us conclude by constructing the forgetful functor

$$\text{Mod}^{\mathcal{O}}(\mathcal{C})^{\otimes} \rightarrow \mathcal{C}^{\otimes}.$$

**Exercise 3.11.** Show that any isomorphism in  $\mathcal{O}^{\otimes}$  is semi-inert. In particular, there is a functor

$$\Delta : \mathcal{O}^{\otimes} \rightarrow \mathcal{K}_{\mathcal{O}}, X \mapsto \text{id}_X.$$

**Exercise 3.12.** Show that the diagram

$$\begin{array}{ccc}
 & \mathcal{O}^\otimes & \\
 & \swarrow \parallel & \searrow \parallel \\
 \mathcal{O}^\otimes & \xleftarrow{\text{ev}_0} & \mathcal{K}_0 & \xrightarrow{\text{ev}_1} & \mathcal{O}^\otimes \\
 & & \downarrow & & \\
 & & & & 
 \end{array}$$

induces a natural functor

$$\text{ev}_{0,*} \circ \text{ev}_1^*(\mathcal{C}^\otimes) \rightarrow \mathcal{C}^\otimes$$

defined over  $\mathcal{O}^\otimes$ . Its restriction

$$\text{Mod}^0(\mathcal{C})^\otimes \rightarrow \mathcal{C}^\otimes$$

is called the **forgetful functor**.

#### APPENDIX A. FLAT FIBRATIONS

A.1. Let  $q : X \rightarrow Y$  be a categorical fibration between quasi-categories. As mentioned in Remark 2.13, the functor  $q_* : (\text{Cat}_\infty)_{/X} \rightarrow (\text{Cat}_\infty)_{/Y}$  exists iff  $q$  is a *flat* categorical fibration. The latter condition has a nice characterization via simplicial sets, which says  $q$  is flat iff for any  $\Delta^2 \rightarrow Y$ , the map  $X \times_Y \Lambda_1^2 \rightarrow X \times_Y \Delta^2$  is a categorical equivalence. Informally speaking, this means for any chain  $u \xrightarrow{f} v \xrightarrow{g} w$  in  $Y$ , the *correspondence* between  $X_u$  and  $X_w$  over  $g \circ f$  is the composition of the correspondences over  $f$  and  $g$ .

A.2. **Suggested readings.** HA.B.3.